MEMORY EQUIPMENT OPTIONS

The basic 4K or 8K memory supplied with the PDP-8/E may be expanded in increments of 4K or 8K. One memory extension and timeshare control is required whenever the system is expanded above 4K. Expansion in 4K increments is achieved by adding one MC8-E Core Memory with Memory Extension and Time Share Control to the 4K system, and then adding up to 6 units of the MM8-E 4K Core Memory. The 4K system is expanded in 8K increments with the addition of one MC8-EJ 8K Core Memory with Memory Extension and Timeshare Control, plus one or two units of MM8-EJ 8K Core Memory. The 8K memory extensions and 4K memory extensions may be mixed in one processor. In any event, either one MC8-E, one MC8-EJ or one KM8-E (described below) is required whenever the basic 4K memory is expanded. The memory extension and timeshare control portion of the MC8-E and MC8-EJ are programmed in the same manner as the KM8-E.

KM8-E Memory Extension and Time-Share Option

This option provides the user with two primary capabilities. The memory extension portion extends the addressing capabilities of the machine from 4069 words up to 32,768 words. The time-share portion enables the computer to operate in either the normal manner (Executive Mode) or the User Mode. User Mode enables the machine to function in a time-sharing environment in which a user program is prevented from disturbing or interfering with another user program. The KM8-E option is packaged on one PDP-8/E module that plugs into the OMNIBUS. This option is required whenever memory capacity is extended beyond 4096 words.

The functional circuit elements which make up the memory extension control perform as follows:

Instruction Field Register (IF)—The IF is a three-bit register that serves as an extension of the PC. The contents of the IF determine the field from which all instructions are taken and the field from which operands are taken in directly-addressed AND, TAD, ISZ, or DCA instructions. Depressing the console EXTD ADDR LOAD switch transfers the instruction field in SWITCH REGISTER bits 6 through 8 into the IF register. During a JMP or JMS instruction, the IF is set by transfer of information from the instruction buffer register. When a program interrupt occurs, the contents of the IF are automatically stored in bits 0 through 2 of the save field register for restoration to the IF from the instruction buffer register at the conclusion of the program interrupt subroutine.

Data Field Register (DF)—This three-bit register determines the memory field from which operands are taken in indirectly-addressed AND, TAD, ISZ, or DCA instructions. Depressing the console EXTD ADDR LOAD switch transfers the SWITCH REGISTER bits 9 through 11 into the DF register. During a CDF instruction, the DF register is loaded from MD6-8 to establish a new data field. When a program interrupt occurs, the contents of the DF are automatically stored in bits 3-5 of the save field register. The DF is set by a transfer of information from save field register bits 3 through 5 by the RMF instruction. This action is required to restore the data field at the conclusion of the program interrupt subroutine.

Instruction Buffer Register (IB)—The IB serves as a three-bit input buffer for the instruction field register. All field number transfers into the instruction field register are made through the instruction buffer, except transfers from the operator's console switches. The IB is set by depressing of the console EXTD ADDR LOAD switch in the same manner as the instruction field register. A CIF microinstruction loads the IB with the programmed field on MD6-8. An RMF microinstruction transfers save field register bits 0 through 2 into the IB to restore the instruction field that existed before a program interrupt.

Save Field Register (SF)—When a program interrupt occurs, this seven-bit register is loaded from the user build flip-flop, and the IR and UF registers. The SF register is loaded during the cycle in which the program count is stored at address 0000 of the JMS instruction forced by a program interrupt request, then the instruction field and data field registers are cleared. An RMF instruction can be given immediately before exit from the program interrupt subroutine to restore the instruction field and data field by transferring the SF into the IB and the DF registers. (Also, see GTF and RTF instructions.)

Extended Address Gating—This logic consists of an output gating structure and control logic for gating the extended memory field address to core memory. The contents of the IF register are placed on the EMA0-2 lines unless an AND I, TAD I, ISZ I, or DCA I instruction is encountered. If such an instruction is encountered, the contents of the IF are placed on EMA 0-2 for the Fetch and Defer cycles, and the contents of the DF are placed on EMA0-2 for the Execute cycle. The extended memory field address is changed only at TP4 and remains available for the entire memory cycle.

Data Transfer Gating—This gating allows the contents of the save field register, instruction field register, or the data field register to be strobed into the accumulator via DATA lines 6-11. During an RIB or GTF instruction, bits 6 through 11 of the AC receive contents of the save field register. During an RIF instruction, bits 6 through 8 of the AC receive the contents of the instruction field register. During an RDF instruction, bits 6 through 8 of the AC receive the contents of the data field register.

Device Selector and Instruction Decoding—Bits 3 through 5 of the IOT instruction are decoded to produce the IOT command pulses for the memory extension control. Bits 6 through 8 of the instruction are not used for device selection since they specify a field number in some commands. Therefore, the select code for this device selector is designated as 2N. Bits 9 through 11 are also decoded to implement specific commands. The instruction decoding logic is common to the time-share portion of the KM8-E option.

Programming

Instructions associated with the extended memory portion KM8-E option are defined below:

Get Flags (GTF)

Octal Code: 6004
Operation: Reads the contents of the interrupt inhibit flip-flop, and the SF register to AC3, AC5-11 respectively. The other AC bits are loaded with information from the CPU and the EAE; i.e., link, greater-than-flag, interrupt bus, interrupt on.

Restore Flags (RTF)

Octal Code: 6005
Operation: Loads the user buffer flip-flop, the instruction buffer register, and the data field register with the contents of AC bits 5, 6-8, and 9-11 and inhibits processor interrupts until the next JMP or JMS instruction. At the conclusion of the JMP or JMS instruction, the contents of the user buffer flip-flop and the instruction buffer register are transferred into the user field flip-flop and the instruction field register, respectively. The contents of the other AC bits are loaded into the CPU and EAE to cause the converse of the GTF instruction. The Interrupt On flip-flop in the CPU is unconditionally set by this instruction.

Change to Data Field N (CDF)

Octal Code: 62N1
Operation: Loads the data field register with the program-selected field number (N = 0 to 7). All subsequent memory requests for operands are automatically switched to that data field, except for directly-addressed AND, TAD, ISZ, or DCA instructions.

Change to Instruction Field N (CIF)

Octal Code: 62N2
Operation: Loads the instruction buffer register with the program-selected field number (N = 0 to 7) and inhibits processor interrupts until the next JMP or JMS instruction. At the conclusion of either of these instructions, the contents of the instruction buffer register are transferred into the instruction field register.

Change Data Field, Change Instruction Field (CDF, CIF)

Octal Code: 62N3
Operation: Performs the combination of CDF and CIF operations.

Read Data Field (RDF)

Octal Code: 6214
Operation: ORs the contents of the data field register into bits 6-8 of the AC. All other bits of the AC are unaffected.

Read Instruction Field (RIF)

Octal Code: 6224
Operation: ORs the contents of the instruction field register into bits 6-8 of the AC. All other bits of the AC are unaffected.

Read Interrupt Buffer (RIB)

Octal Code: 6234
Operation: ORs the contents of the save field register (which is loaded from the instruction and data field during a program interrupt) into bits 6-8 and 9-11 of the AC, respectively. Thus, AC 6-11 contains the instruction and data fields that were in use before the last program interrupt. AC 5 is loaded by the time-share bit of the save field register. All other bits of the AC are unaffected.

Restore Memory Field (RMF)

Octal Code: 6244
Operation: Restores the contents of the save field register (which is loaded from the instruction and data field during a program interrupt) into the instruction buffer, the data field register, and the user buffer (if time share option is enabled). This command is used upon exit from the program interrupt subroutine in another field.

Instructions and data are accessed from the currently assigned instruction and data fields, where instructions and data may be stored in the same or different memory fields. When indirect memory references are executed, the operand address refers first to the instruction field to obtain an effective address, which, in turn, refers to a location in the currently assigned data field. All instructions and operands are obtained from the field designated by the contents of the instruction field register, except for indirectly addressed operands, which are specified by the contents of the data field register. In other words, the DF is effective only in the execute cycle that is directly preceded by the defer cycle of a memory reference instruction, as follows:

Indirect (Bit 3)	Page or Z Bit (Bit 4)	Field In IF	Field In DF	Effective Address
0	0	m	n	The operand is in page 0 of field m at the page address specified by bits 5 through 11.
0	1	m	n	The operand is in the current page of field m at the page address specified by bits 5 through 11.
1	0	m	n	The absolute address of the operand in field n is taken from the contents of field m located in page 0 designated by bits 5 through 11.

1	1	m	n	The absolute address of the operand in field n is taken from the contents of field m located in the current page, designated by bits 5 through 11.

Each field of extended memory contains eight auto-index registers in addresses 10 through 17. For example, assume that a program in field 2 is running (IF = 2) and using operands in field 1 (DF = 1) when the instruction TAD I 10 is fetched. The defer cycle is entered (bit 3 = 1), and the contents of location 10 in field 2 are read, incremented, and rewritten. If address 10 in field 2 originally contained 4321, it now contains 4322. In the execute cycle, the operand is fetched from location 4322 of field 1. Program control is transferred between memory fields by the CIF instruction. The instruction does not change the instruction field directly, as this would make it impossible to execute the next sequential instruction; instead, it loads the new instruction field in the IB for automatic transfer into the IF when either a JMP or JMS instruction is executed. The DF is unaffected by the JMP and JMS instructions.

The 12-bit program counter is set in the normal manner and, because the IF is an extension on the most significant end of the PC, the program sequence resumes in the new memory field following a JMP or JMS. Entry into a program interrupt is inhibited after the CIF instruction until a JMP or JMS is executed.

NOTE

The IF is not incremented if the PC goes from 7777 to 0000. This feature protects the user from accidentally entering a nonexistent field.

To call a subroutine that is out of the current field, the data field register is set to indicate the field of the calling JMS, which establishes the location of the operands as well as the identity of the return field. The instruction field is set to the field of the starting address of the subroutine. The following sequence returns program control to the main program from a subroutine that is out of the current field.

```
/PROGRAM OPERATIONS IN MEMORY FIELD 2
/INSTRUCTION FIELD = 2; DATA FIELD = 2
/CALL A SUBROUTINE IN MEMORY FIELD 1
/INDICATE CALLING FIELD LOCATION BY THE CONTENTS OF THE DATA
 FIELD

         CIF 10          /CHANGE TO INSTRUCTION
                         /FIELD 1 = 6212
         JMS I SUBRP     /SUBRP = ENTRY ADDRESS
         CDF 20          /RESTORE DATA FIELD
           .
           .
           .
SUBRP,   SUBR            /POINTER
                         /CALLED SUBROUTINE, LOCATED IN
                         /FIELD 1
```

```
SUBR,   0               /RETURN ADDRESS STORED HERE
        CLA
        RDF             /READ DATA FIELD INTO AC
        TAD RETURN      /CONTENTS OF THE AC = 6202 +
                        /DATA FIELD BITS
        DCA EXIT        /STORE INSTRUCTION SUBROUTINE
          .             /NOW CHANGE DATA FIELD IF DESIRED
          .
          .
          .
EXIT,                   /A CIF INSTRUCTION
        JMP I SUBR      /RETURN TO CALLING PROGRAM
RETURN, CIF             /USED TO CALCULATE EXIT
                        /INSTRUCTION
```

When a program interrupt occurs, the current instruction and data field numbers are automatically stored in the 6-bit save field register, then the IF and DF are cleared. The 12-bit program count is stored in location 0000 of field 0 and program control advances to location 0001 of field 0. At the end of the program interrupt subroutine, the RMF instruction restores the IF and DF from the contents of the SF. Alternatively, the GTF and RTF instructions may be used to handle the Save Field and Link information. The following instruction sequence at the end of the program interrupt subroutine continues the interrupted program after the interrupt has been processed:

```
          .             /RESTORE MQ IF REQUIRED
          .
          .
          .             /RESTORE L IF REQUIRED
          .
          .
          .
          .
        CLA
        TAD AC          /RESTORE AC
        RMF             /LOAD IB FROM SF
        ION             /TURN ON INTERRUPT SYSTEM
        JMP I 0         /RESTORE PC WITH CONTENTS OF
                        /LOCATION 0 AND LOAD IF FROM IB
OR
0,      0               /PC STORAGE
        DCA ACSV        /SAVE AC,
        MQA CLA
        DCA MQSV        /MQ,
        GTF
        DCA FLAGS
          .
          .
          .
        CLA
        TAD MQSV
```

```
MQL             /RESTORE MQ
TAD FLAGS
RTF             /REPLACE FLAGS, ION
CLA
TAD ACSV        /AC
JMP I 0         /AND EXIT
```

Time-Share Description

The Time-Share portion of the KM8-E operates in two modes as denoted by the user flag (UF) flip-flop. When the UF flip-flop is in the logic 1 state, operation is in the user mode and a user program is running in the central processor. When the UF flip-flop is in the logic 0 state, operation is in the executive mode and the time-sharing system's monitor is in control of the central processor. The four instructions (CINT, SINT, CUF, and SUF) are used by the time-sharing system's monitor in the executive mode and are never used by a user program. If a user program attempted to use one of these instructions, execution of the instruction would be blocked (see next paragraph). The KM8-E option adds the necessary hardware to the PDP-8/E to implement these instructions.

In executive mode, the computer operates normally. When the computer is operated in user mode, operation is normal except for IOT, HLT, LAS, and OSR instructions. When one of these instructions is encountered, the hardware inhibits the normal instruction sequence (other than rewriting the instruction in memory), and generates an interrupt at the end of the current memory cycle by setting the UINT flip-flop. The time-sharing system's monitor program then analyzes the source of interrupt, and takes appropriate action.

The time-share option requires at least 8K of core memory; thus, it is packaged with the memory extension option. A jumper on the KM8-E module is used to select the time-share function. The module is shipped with this jumper in place (time-share function disabled).

Programming

Instructions associated with the time-share portion of the KM8-E are defined as follows:

Clear User Interrupt (CINT)

Octal Code: 6204
Operation: Clears the user interrupt flip-flop.

Skip on User Interrupt (SINT)

Octal Code: 6254
Operation: Increments the PC when the user interrupt flip-flop is set so the next sequential instruction is skipped.

Clear User Flag (CUF)

Octal Code: 6264
Operation: Clears the user buffer flip-flop.

NOTE

If the machine is stopped while in user mode, the user flag (UF) is cleared by operating the extended address load key (EXT ADDR LOAD).

Octal Code: 6274

Operation: Sets user buffer flip-flop and inhibits processor interrupts until the next JMP or JMS instruction. At the conclusion of either of these instructions, the content of the user buffer flip-flop is transferred into the user field flip-flop.

MP8-E Memory Parity

The memory parity option adds the circuits required to generate, store, and check the parity of memory words. This option replaces the 12-bit memory system with, effectively, a 13-bit system by adding the generating and storage capabilities for the parity bit. Odd parity (odd number of binary ones in the 13-bit word) is generated and stored for each word entered into memory. Parity is formed for each word retrieved from memory and this result is checked against its stored parity bit. If the two differ, a parity error flag is set to indicate that an error occurred. This flag is normally connected to the program interrupt system to cause the computer to enter a program interrupt subroutine for locating the interrupt source. Once the interrupting source is located, the computer enters an appropriate service routine to service the error condition. This routine can repeat the program step in which the error occurred to verify the error condition, can perform a simple read/write check for the error's address, or can determine machine status for the error detected and re-establish or print out these conditions, and then halt. The routine can also return the machine to the main program.

The MP8-E option consists of three PDP-8/E modules that plug into the OMNIBUS. Two of these modules (X-Y Driver and Current Source, and Core Stack) are identical to those of the MM8-E basic core memory and use the same addressing methods. However, only eight bits of the possible 12 bits are used. These eight-bit locations correspond to the eight possible memory fields and store up to 32,768 (8 $\times$ 4096) parity bits. The third module (Sense-Inhibit) contains device and operation decoding circuits, field decoding circuits, eight sense amplifiers, an eight-bit register, eight inhibit drivers and circuits for controlling the operations. This module also contains three control and status flip-flops that are controlled by IOT instructions. These flip-flops select odd or even parity generation and checking, enable or disable interrupts for parity errors, and store a parity error condition.

The following routine initializes the parity bits for a read-only or write-protected memory:

```
                        /INITIALIZE LOC 10 WITH STARTING ADD.
                        /TURN OFF PARITY INTERRUPT
                        /SET COUNTER
LOOP,   TAD I 10        /READ DATA, REWRITE PARITY
        ISZ COUNT
        JMP LOOP        /CONTINUE UNTIL DONE
                        /CLEAR PARITY ERROR FLAG
                        /TURN ON PARITY INTERRUPT
```